

The parallel intentionally fully abstract games model of PCF

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Abstract—We describe a framework for truly concurrent game semantics of programming languages, based on Rideau and Winskel’s concurrent games on event structures. The model supports a notion of innocent strategies that permits concurrent and non-deterministic behaviour, but which coincides with traditional Hyland-Ong innocent strategies if one restricts to the deterministic sequential case. In this framework we give an alternative interpretation of Plotkin’s PCF, that takes advantage of the concurrent nature of strategies and formalizes the idea that although PCF is a sequential language, certain sub-computations are independent and can be computed in a parallel fashion. We show that just as Hyland and Ong’s sequential interpretation of PCF, our parallel interpretation yields a model that is intentionally fully abstract for PCF.

I. INTRODUCTION

Regardless of mathematical elegance, *partial order* models of concurrent computation are in principle more informative than their interleaving counterparts: they avoid the state explosion problem inherent to interleavings, and retain explicit information on causality. This can be useful for instance for the purposes of error diagnostics, or security analysis. However, although we have truly concurrent models for simple process languages such as CCS, extracting partial order models from source code remains a challenge, especially if one considers rich concurrent programming languages with complex computational features such as higher-order, state or exceptions.

In order to construct compositionally a fine-grained representation of the execution of higher-order programs, *game semantics* is a powerful tool. Game semantics proposes to see computation as an interaction between agents (strategies) exchanging messages, hence reducing higher-order computation to the exchange of first-order tokens. Thanks to this methodology, game semantics has not only given intensionally fully abstract models of PCF [10, 1] but also pushed beyond the purely functional setting and given effectively presentable fully abstract models of higher-order programming languages with rich computational features such as control or state.

Most games models for concurrent programming languages [7, 12], however, are based on interleavings. Several truly concurrent frameworks for game semantics have been proposed [6, 14, 16, 9], but have yet to be applied to the semantics of programming languages beyond CCS or linear logic – this is in part due to the fact that truly concurrent notions of strategies are combinatorially more complex than their interleaved counterparts, and are more subtle to handle. Moreover, changing

the basic metalanguage for game semantics means losing a whole body of work, such as the pleasing characterisation offered by Hyland-Ong (HO) games of various computational effects in terms of conditions on sequential strategies. In this paper, we aim to fill in this gap. Our first contribution is a framework for truly concurrent game semantics, that comprises representations of computational features such as non-determinism and concurrency while containing as a sub-case the usual HO games – in particular we have notions of visibility and innocence, that in the sequential deterministic case coincide with standard HO innocence. We believe that this framework should prove adequate for further developments on truly concurrent games models of programming languages.

Our second contribution is an application of our framework to give a *parallel* intensionally fully abstract model for PCF. That might seem counter-intuitive, given the status of PCF as a paradigmatic *sequential* language. However, although all the primitives of PCF are purely sequential, that does not mean that its *implementation* has to be sequential. In particular, operations such as the conditional $\text{if} : \mathbb{B} \rightarrow \mathbb{B} \rightarrow \mathbb{B} \rightarrow \mathbb{B}$ could in principle be optimized by evaluating all three arguments in parallel, and returning the adequate one according to the result of the first argument. Standard game semantics of PCF specify explicitly an evaluation order, and forbids this operational reading of *if*. Our game semantics, while authorizing the execution as prescribed by the sequential game semantics of PCF, will make this parallel computation official and express this parallelism in a truly concurrent manner. Despite this added intensional behaviour, our conditions will be enough to guarantee that our strategies are extensional, and that their extensional collapse is the fully abstract model for PCF.

Related works. On the game-theoretic front, the present contribution was made possible by the recent developments [16, 4, 21] initiated by Rideau and Winskel around a framework for game semantics based on event structures – our basic setting is an extension of the games with symmetry of [4]. The framework of [16] generalizes earlier approaches to deterministic truly concurrent notions of games [2, 14, 6], used in particular by Melliès to build a fully complete model of full propositional linear logic [13]. Also relevant is the work of Hirschowitz and Pous [9, 8], which gives a fully abstract model of CCS *w.r.t.* fair testing, based on a notion of strategies as presheaves on an abstract category of plays.

Outline. In Section II we will present our notions of games

Types. $A, B ::= \mathbb{B} \mid \mathbb{N} \mid A \rightarrow B$

Terms. $M, N ::= x \mid \lambda x. M \mid M N \mid Y \mid$
 $\text{tt} \mid \text{ff} \mid \text{if } M_1 M_2 M_3 \mid$
 $n \mid \text{succ } M \mid \text{pred } M \mid \text{iszero } M$

Fig. 1. The language PCF

and strategies and compare them with the HO innocent well-bracketed strategies. In Section III, we present the mathematical foundations for our model, and describe a sound interpretation of PCF. Finally, in Section IV we prove a finite definability result and deduce full abstraction.

II. PARALLEL PCF-STRATEGIES

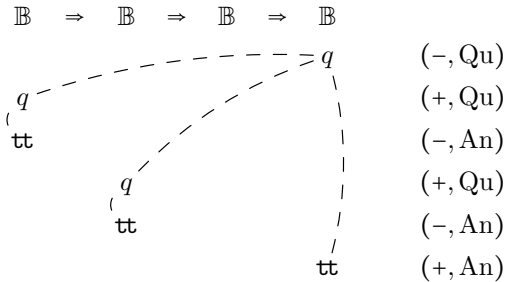
We introduce in Figure 1 the syntax of PCF as used in this paper. In this paper \mathbb{X} will range over **ground types**, *ie.* \mathbb{B} or \mathbb{N} . The typing rules we consider are standard, except for the typing rule for if that is replaced with:

$$\frac{\Gamma \vdash M : \mathbb{B} \quad \Gamma \vdash N_1 : \mathbb{X} \quad \Gamma \vdash N_2 : \mathbb{X}}{\Gamma \vdash \text{if } M N_1 N_2 : \mathbb{X}}$$

Note that if N_1, N_2 have arbitrary type $A = A_1 \rightarrow \dots \rightarrow A_n \rightarrow \mathbb{X}$, if $M N_1 N_2$ can still be used and is considered syntactic sugar for $\lambda x_1 \dots x_n. \text{if } M (N_1 x_1 \dots x_n) (N_2 x_1 \dots x_n)$. Terms of PCF are executed following the standard call-by-name big-step operational semantics, yielding an evaluation relation \Downarrow between closed terms and **values** (*ie.* constants of ground type or λ -abstraction). As usual we write $M \Downarrow$ if $M \Downarrow v$ for some v . In this section, we aim to explain operationally what our model computes, and how it represents programs.

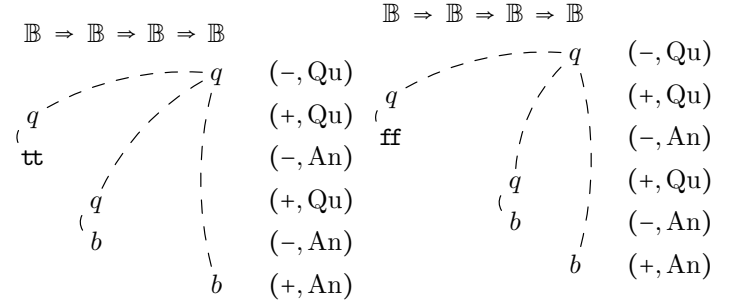
A. Sequential PCF-strategies presented operationally

HO game semantics formalizes the intuition that a program is a strategy having a dialogue with its execution environment. A possible dialogue on the type $\mathbb{B} \Rightarrow \mathbb{B} \Rightarrow \mathbb{B} \Rightarrow \mathbb{B}$ could be:



Each move is either Player or Opponent, and is either a Question or an Answer. Questions correspond to variable calls, whereas Answers indicate a call terminating. The dashed lines between moves (traditionally called *justification pointers*) convey information about thread indexing; in this example they are redundant but become required at higher types. **Sequential innocent strategies** consist of sets of dialogues as above, where Opponent moves are *justified* by the preceding one

– such dialogues are known as *P-views*. A strategy for a PCF term contains several such dialogues, specifying the term entirely. In our example, the dialogue is actually a branch of the strategy for if, that interrogates its left argument first, then the second or third depending on the result. This strategy contains in total four maximal P-views, described by:

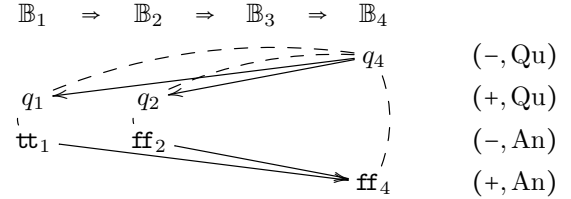


Such non-empty sets of P-views (satisfying further conditions called determinism and well-bracketing) are called **sequential PCF-strategies**, or PCF-strategies for short.

B. PCF_{||}-strategies

Instead of investigating its arguments sequentially, an optimized implementation of if could evaluate the second and third arguments before they are needed, without waiting for the first call to terminate. This intuition leads to a different notion of strategy, and a different interpretation of PCF.

1) *A partial order for if*: States of this optimized strategy for if will no longer be total orders but *partial* orders. For instance, a state of if where the first argument returns tt could be (labelling the distinct copies of \mathbb{B} for clarity):

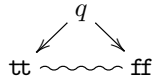


Additionally to the justification pointers, the strategy is now equipped with a relation \rightarrow indicating *immediate causality* (omitted for readability when it coincides with vertical juxtaposition). In the sequential diagrams of Subsection II-A these were not required since in a sequential setting immediate causality is chronological contiguity. In the diagram above, any two events not related by the transitive closure of immediate causality are concurrent, and occur independently. So (amongst other actions) if interrogates its two first arguments in parallel, and is able to answer ff if the first argument returns tt and the second ff. Instead of being a total order, a “branch” of a strategy will have to be generalized to a partial order – of course “branch” here is misused, instead we will speak of a *prime*, *ie.* a partial order with a maximum element. As we will see soon, the full strategy will contain other primes for different executions leading to an answer in \mathbb{B}_4 .

2) *Event structures*: Even though our strategies for PCF are all deterministic, we aim for a framework that can accommodate non-deterministic behaviour. Accordingly, we will use *event structures* [18]: a concurrent analogue of trees, providing a description of systems that features both independence of events, and non-determinism. Formally, an **event structure** is (E, \leq, Con) with E a set of events, \leq a partial order indicating *causal dependency*, and Con a nonempty *consistency relation* consisting of finite subsets of E , such that:

$$\begin{aligned} \{e' \in E \mid e' \leq e\} &\text{ is finite for all } e \in E \\ \{e\} &\in \text{Con for all } e \in E \\ Y \subseteq X \in \text{Con} &\implies Y \in \text{Con} \\ X \in \text{Con} \ \& \ e \leq e' \in X &\implies X \cup \{e\} \in \text{Con} \end{aligned}$$

For example, the event structure for the non-deterministic strategy returning a random boolean is:



where the wiggly line indicates *immediate conflict*, i.e. the consistent sets are those not containing both tt and ff .

Event structures will sometimes be equipped with a **polarity** function $\text{pol} : E \rightarrow \{-, +\}$ specifying whether an event e is Player ($\text{pol}(e) = +$) or Opponent ($\text{pol}(e) = -$). An event structure E with pol is called an **event structure with polarities**, or *esp* for short.

3) *Notations*: For $e, e' \in E$, we write $e \rightarrow e'$ for **immediate causality**, i.e. $e < e'$ and for all e'' such that $e \leq e'' \leq e'$, we have $e = e''$ or $e'' = e'$. If E is an event structure, we write $\mathcal{C}(E)$ for the set of **configurations** of E , comprising the finite subsets $x \subseteq_f E$ such that $x \in \text{Con}$ and x is down-closed, i.e. for all $e \in x$, $e' \leq e$, we also have $e' \in x$. For $e \in E$, we write $[e]$ the corresponding *prime configuration* or *prime* for short, defined as $[e] = \{e' \in E \mid e' \leq e\}$. Finally if E has polarity, we write $x \subseteq^+ y$ (resp. $x \subseteq^- y$) if $x \subseteq y$ and $\text{pol}(y \setminus x) \subseteq \{+\}$ (resp. $\text{pol}(y \setminus x) \subseteq \{-\}$).

C. Concurrent pre-strategies playing on arenas

We now describe more formally our notion of strategies.

1) *Arenas*: As usual in HO game semantics, the types of PCF will be interpreted as *arenas*.

Definition II.1. An **arena** is a tuple $(A, \leq, \text{pol}, \lambda)$ such that $(A, \leq, \mathcal{P}_f(A), \text{pol})$ is a countable *esp* satisfying:

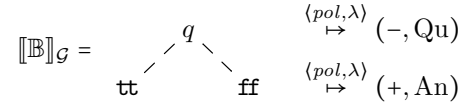
$$\begin{aligned} \forall a, a', a'' \in A, a \leq a'' \ \& \ a' \leq a'' \implies a \leq a' \vee a' \leq a \\ \forall a, a' \in A, a \rightarrow a' &\implies \text{pol}(a) \neq \text{pol}(a') \\ \forall a \in A, a \in \min(A) &\implies \text{pol}(a) = - \end{aligned}$$

with $\min(A)$ the set of minimal events of A and $\lambda : A \rightarrow \{\text{Qu}, \text{An}\}$ a *Question/Answer labeling function*, such that:

$$\begin{aligned} \forall a \in A, a \in \min(A) &\implies \lambda(a) = \text{Qu} \\ \forall a_1, a_2 \in A, a_1 \rightarrow a_2 &\implies \lambda(a_1) = \text{Qu} \end{aligned}$$

Readers familiar with HO games will recognize the notion of arenas of [10], with the partial order \leq primitive rather than immediate causality \rightarrow (traditionally written \vdash in HO games).

Types of PCF_{\parallel} are interpreted as arenas. In particular, the basic type \mathbb{B} is interpreted as (read from top to bottom):



The type \mathbb{N} for natural numbers is interpreted similarly, but with a countably infinite number of answers $0, 1, \dots$. In the remainder of this paper we will often omit the semantic brackets and have the same notations for constructions on types and arenas, except when brackets are useful for disambiguation.

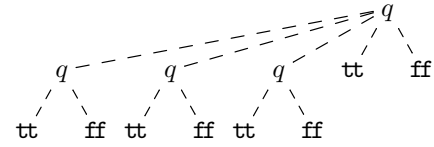
Given two arenas A and B , $A \Rightarrow B$ is defined as having:

- Events, $\{(1, (b, a)) \mid a \in A \ \& \ b \in \min(B)\} \cup \{(2, b) \mid b \in B\}$.
- Causality,

$$\begin{aligned} \{((1, (b, a_1)), (1, (b, a_2))) \mid a_1 \leq a_2 \ \& \ b \in \min(B)\} \cup \\ \{((2, b_1), (2, b_2)) \mid b_1 \leq b_2\} \cup \\ \{((2, b), (1, (b, a))) \mid a \in A \ \& \ b \in \min(B)\} \end{aligned}$$

- Polarity, $\text{pol}((1, (b, a))) = -\text{pol}(a)$, $\text{pol}((2, b)) = \text{pol}(b)$.
- Qu/An labeling, $\lambda((1, (b, a))) = \lambda(a)$ and $\lambda((2, b)) = \lambda(b)$.

The reader familiar with sequential HO games will recognize here the usual arrow arena construction $A \Rightarrow B$. With these definitions, the arena $\mathbb{B} \Rightarrow \mathbb{B} \Rightarrow \mathbb{B} \Rightarrow \mathbb{B}$ of if is:



2) *Expanded game*: In sequential HO games, strategies do not play on arenas but rather on the derived game of *plays with pointers* [10], where the same move can be reused at will. We perform a similar construction here: from each arena we derive a game where events can be played as many times as required. Unlike sequential HO games though, rather than building a tree of plays we build an event structure where duplicated events retain the causal structure of the arena. Unlike sequential plays where distinct copies of the same move are kept apart chronologically, here we handle explicit copy indices.

To define it we need the notion of an *indexing function*. If $a \in A$ for A an arena, an **indexing function** for a is:

$$\alpha : [a] \rightarrow \omega$$

associating a copy index (a natural number) to each dependency of a , itself included. For $\alpha : [a] \rightarrow \omega$, we write $\text{lbl } \alpha = a$ for its **label** and $\text{ind } \alpha = \alpha(\text{lbl } \alpha)$ for its **index**, i.e. the copy index it associates to the maximum of its domain.

Definition II.2. Let A be an arena. There is an *esp* $!A$ having:

- Events: indexing functions,
- Causality: $\alpha \leq \alpha'$ defined as $\text{lbl } \alpha \leq \text{lbl } \alpha'$ and α, α' agree on their common domain,
- Consistency: trivial – all finite sets consistent.
- Polarity: $\text{pol}(\alpha) = \text{pol}(\text{lbl } \alpha)$.

Events $\alpha \in !A$ inherit a Question/Answer labeling from A . We also define $!^+A$ duplicating only positive events, i.e. with $\alpha : [a] \rightarrow \omega$ such that for all $a' \leq a$ with $\text{pol}(a') = -$, $\alpha(a') = 0$.

Configurations of $!A$ correspond to Boudes' *thick subtrees* [3] of arenas: a configuration $x \in \mathcal{C}(!A)$ visits a prefix of the arena A , but with branches duplicated at will. For now $!A$ is still technically an arena, but we will not consider it as such: in the compositional development to come we will equip it with an additional notion of *symmetry* (see Section III). Symmetry is key for the compositional structure of our model, but we can ignore it for the purposes of this section.

3) *Pre-strategies*: A **map** from event structure A to event structure B can be thought of as a simulation of A within B . Formally it is a function $f : A \rightarrow B$ on events, which:

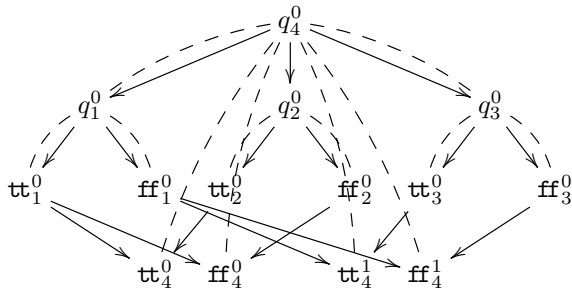
- Preserves configurations: for all $x \in \mathcal{C}(A)$, $f x \in \mathcal{C}(B)$,
- Is locally injective: for all $x \in \mathcal{C}(A)$, for all $e, e' \in x$, $f e = f e' \implies e = e'$.

Additionally if A and B have polarity, f preserves it.

Now, a **pre-strategy** on arena A is defined as a map of esp.:

$$\sigma : S \rightarrow !A$$

This definition allows us to make formal the diagrams of the previous subsection: the nodes represent events of S , annotated with the label of the corresponding event in the game (as obtained through σ). The arrow \rightarrow represents immediate causality in S , while the dashed lines $-$ represent the relation induced on S by immediate causality in $!A$. The diagrams of Subsection II-B were incomplete since they omitted copy indices. With these, the full (reduced – see below) pre-strategy for $\text{if} : \mathbb{B}_1 \Rightarrow \mathbb{B}_2 \Rightarrow \mathbb{B}_3 \Rightarrow \mathbb{B}_4$ will have:



The superscript of a move s indicates its copy index, i.e. $\text{ind}(\sigma s)$. The full index function can be recovered from the copy index annotations of the dependencies in $!A$, as expressed by the dashed lines. As there are several compatible ways of returning tt (resp. ff), the pre-strategy has to use distinct copy indices for the projection to the game to be locally injective.

Any sequential innocent strategy on an arena A in the sense of HO games [10] can be represented as a pre-strategy $\sigma : S \rightarrow !A$, where S is the forest of correct P-views, and a P-view is sent by σ to its latest move along with a copy index, chosen as to avoid collisions. However, there are many more pre-strategies, some non-deterministic or concurrent.

D. Conditions for PCF_{\parallel} -strategies

Of course, not all pre-strategies $\sigma : S \rightarrow !A$ or relevant for PCF. For the purposes of this paper, we need to describe a class of strategies containing our strategy for if and closed under composition (as will be made formal in the next section), but still small enough so that strategies within it have the same distinguishing power as terms of PCF.

A first remark is that PCF_{\parallel} -strategies should be uniform, in the sense that their behaviour should not depend on Opponent's choice of copy indices. This seemingly elementary notion actually requires some machinery to formalize in a compositional setting. However, since in this section we only examine pre-strategies as purely static objects, we will content ourselves in examining **reduced pre-strategies**, which – as in the examples above – only acknowledge Opponent moves of copy index 0, i.e. play on the expanded game $!^+A$. Uniformity aside, we now examine which properties PCF_{\parallel} -strategies should satisfy w.r.t. causality and consistency. We introduce the conditions from the more robust to the more PCF-specific.

1) *Courteous and receptive pre-strategies*: Following [16] we define a **strategy** as a pre-strategy $\sigma : S \rightarrow !^+A$ which is:

- **Receptive**: For all $x \in \mathcal{C}(S)$, for all $\sigma x \sqsubseteq^- y'$, there is a unique $x \subseteq x' \in \mathcal{C}(S)$ such that $\sigma x' = y'$.
- **Courteous**: For all $s_1, s_2 \in S$ such that $s_1 \rightarrow s_2$, if $\text{pol}(s_1) = +$ or if $\text{pol}(s_2) = -$ then $\sigma s_1 \rightarrow \sigma s_2$.

In game semantics, receptivity is always present in one way or another. It is explicit and named *contingent completeness* in [10], but in most works on game semantics it is hard-wired by asking that strategies contain only plays of even length (Opponent extensions being always present, they bring no additional information).

On the other hand, courtesy expresses that a strategy can only add causal links from negative to positive events. Of course it makes sense that if the rules of the game authorize Opponent to make a move, Player can not force them to wait. What might be more surprising is that Player is not capable of putting additional causal links between their own (positive) moves! One way to understand that is that strategies are thought of as interacting in a distributed fashion, over a network with an uncontrolled latency. So even though Player might want to play their moves in a specific order, they cannot control in what order they will reach their Opponent – this intuition is made formal by the result of [16] that receptive courteous strategies are exactly those that are unchanged by their composition with *copycat*.

Courtesy has a more chaotic history than receptivity: in essence, it is already present in interleaving-based game semantics for concurrency, where it forces strategies to be saturated under a number of permutations, in effect breaking unauthorized causal links. Under this form it dates back to [12], and is called *saturation* in [7] – a name that does not fit here, since it is not a saturation condition. In true concurrency approaches to game semantics, it appears under the name of “innocence” in [16, 6], and “courteousness” in Melliès and Mimram's asynchronous games [14]. We believe innocence is

a slightly misleading name, since (as appears in this paper) it has no direct relationship with Hyland and Ong’s notion of innocence, so we use courtesy instead.

2) *Visibility*: In sequential HO games, strategy branches where Opponent always points to the previous move (*ie.* P -views) correspond to branches of terms. Here this intuition still holds, except that P -views are replaced by *grounded causal chains*. A **grounded causal chain** in S is a sequence of events:

$$\rho_1 \rightarrow \dots \rightarrow \rho_n$$

where $\rho_1 \in \min(S)$. If $\rho_n = s \in S$, we write $\rho \in \text{gcc}(s)$; we also write $|\rho| = n$ for the length of ρ . In our games, grounded causal chains (gccs for short) give a notion of *thread*: if throws three sub-threads, and possible merges some of them later. Each thread should go on independently, until it is merged with another or terminated. In particular, each thread should only use resources introduced within it; we call this *visibility*. A strategy $\sigma : S \rightarrow !^+A$ is **visible** iff:

$$\forall s \in S, \forall \rho \in \text{gcc}(s), \sigma \rho \in \mathcal{C}(!^+A)$$

This amounts to ρ containing all the justifiers (*ie.* immediate dependencies in the game) of Player events. Since σ is courteous, then in $\rho \in \text{gcc}(s)$ causal links from positive to negative events match the causal links of the game – in other words Opponent points to the previous move, so the projection of ρ to the arena A is a P -view.

Visibility is of paramount importance in our development, and is indispensable to the stability under composition of almost all the forthcoming conditions. For visible strategies it makes sense to think of gccs as threads, and most of the remaining conditions restrict how strategies are allowed to generate, merge, or terminate threads.

3) *Innocence*: The *sequential innocent deterministic strategies* (described by a set of P -views) of [10] are the cornerstone of HO games. In our setting, they correspond to the (automatically visible) strategies $\sigma : S \rightarrow !^+A$ additionally satisfying:

- (a) *Backward sequentiality*. The moves available to σ are entirely determined by gccs: for all $s \in S$, $[s]$ is a gcc.
- (b) *Forward sequentiality*. If $s \in S$ and $[s]$ extends by positive distinct s_1, s_2 , then $[s] \cup \{s_1, s_2\} \notin \text{Con}_S$.
- (c) *Determinism*. $\text{Con}_S = \mathcal{P}_f(S)$.

It can be proved that strategies on $!^+A$ satisfying (a), (b) and (c) are in bijection with usual HO innocent strategies on A , up to the choice of copy indices. We note in passing that conditions (a) and (b) together are stable under our forthcoming notion of composition, yielding a notion of non-deterministic sequential innocence, a problem that is known to be difficult in sequential HO games.

However, for our purposes this definition does not fit: our parallel strategy for if fails both (a) and (b). We need to authorize prime configurations $[s]$ to be properly partially ordered, while ensuring that distinct Opponent moves following from the same player move are considered independently by Player. To formalize this, we say that for a strategy $\sigma : S \rightarrow !^+A$ a configuration $x \in \mathcal{C}(S)$ is **normal** if two Opponent moves never share the same justifiers: more formally, for all $s_1, s_2 \in x$

negative events both minimal or such that there is $s \in S$ such that $s \rightarrow s_1, s \rightarrow s_2$ (or equivalently by courtesy, $\sigma s \rightarrow \sigma s_i$), we have $s_1 = s_2$. Then, σ is **innocent** iff its behaviour is specified by normal configurations: for all $s \in S$, $[s]$ is normal.

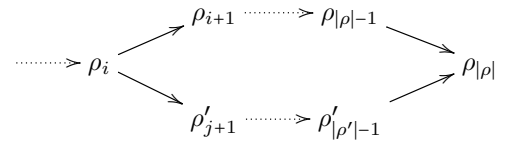
Concurrent innocence no longer implies visibility, although it will be stable under composition only in the presence of visibility. Finally, in the presence of forward sequentiality this definition is equivalent to the one above and so conservatively extends usual sequential innocent strategies.

4) *Well-bracketing*: We first introduce some notation and terminology. If $\sigma : S \rightarrow !^+A$ is a strategy, S inherits a Qu/An labelling from A : an event $s \in S$ is a Question/Answer if $\text{lbl}(\sigma s)$ is. We will sometimes annotate symbols for events with their Qu/An or polarity labelling. We might say for instance “let $s^{-, \text{Qu}} \in S$ ”, meaning that s is an Opponent Question. Given a set $X \in \text{Con}_S$ and $s_1^{\text{Qu}} \in X$, we say that s_1 is **answered** in X if there is an answer $s_2^{\text{An}} \in X$ with $\sigma s_1 \rightarrow \sigma s_2$. Additionally X is **complete** if all the questions of X are answered in X . Finally If ρ is a gcc in S , then we write $\rho_{i \leq j}$ for the segment of ρ lying between indices i and j , endpoints included. We also use the strict variant $\rho_{i < j}$. Finally, ρ_ω denotes the last event of a gcc ρ .

We now give conditions (a), (b) and (c) for well-bracketing. (a) In any gcc ρ in S , if $\sigma \rho_i \rightarrow \sigma \rho_j^{\text{An}}$, then ρ_i is the **pending question** (the latest unanswered question) in $\rho_{< j}$, *ie.* gccs are well-bracketed in the usual sense. Lifting this condition would authorize strategies for control operators such as call/cc [11]. (b) Answering is affine: for $x \in \mathcal{C}(S)$ and two gccs $\rho, \rho' \subseteq x$ such that ρ_ω and ρ'_ω are both answers to the same question, then the greatest common event of $\rho \cap \rho'$ is positive – so ρ and ρ' fork at an Opponent event.

Not having condition (b) means that besides having strategies for the booleans tt and ff one has also, among others, a strategy that answers *both* tt and ff : this strategy creates two copies of its runtime environment and throws two threads, returning tt in the first and ff in the second. There is a striking similarity with the C primitive `fork`, that creates two threads and gives each of them a unique identifier – the details of this connection are left for future work. Before we are done with well-bracketing, there is one last condition to mention:

(c) Take $\rho, \rho' \in \text{gcc}(s^{\text{An}})$ for an answer $s \in S$. Suppose moreover that $\rho_i = \rho'_j$ and that $\rho_{i < |\rho|}$ and $\rho'_{j < |\rho'|}$ are disjoint. Then $\rho_{i < |\rho|}$ and $\rho'_{j < |\rho'|}$ are complete. Graphically, in a diagram as below in S :



the two separate branches are complete – if Player creates two threads, he can only merge them again once all their questions are answered. A similar condition appears in Wall’s games model for syntactic control of interference, see [17].

A visible strategy $\sigma : S \rightarrow !^+A$ is **well-bracketed** when it satisfies all three conditions (a), (b) and (c).

5) *Decomposing PCF_{||}-strategies*: As we shall see later, our interpretation of PCF_{||} will yield strategies satisfying these conditions. Note however that soundness for PCF depends on much weaker conditions (the key one is *single-threadedness*, see Subsection III-D). With the full set of conditions we aim for a much stronger correspondence with syntax; in fact we will describe in Subsection IV-2 a decomposition result akin to the usual definability of sequential innocent strategies.

III. CONCURRENT HYLAND-ONG GAMES

We now describe the mathematical structure required to compositionally generate the strategies described above from PCF terms. The basic ideas behind our mathematical development come from the games on event structures introduced in [16]. To express uniformity of strategies *w.r.t.* copy indices, we use the notion of *symmetry* on event structures [19]. In particular we will use a variant of the games on event structures with symmetry of [4].

A. Event structures with symmetry

A **symmetry** on an event structure (possibly with polarity) E , is usually defined as a *span* $E \xleftarrow{l_E} \tilde{E} \xrightarrow{r_E} E$ in the category of event structures (possibly with polarity), which is an equivalence relation in the categorical sense, l_E and r_E are rigid and are functional bisimulations, expressed as a configuration extension property (see [19]). Intuitively a symmetry on E can be regarded as a *proof-relevant* notion of equivalence relation between configurations of E , in a way that respects the operations available on event structures.

In full generality, symmetry provides an abstract notion of bisimulation in a categorical setting [19]. In the concrete case of event structures, symmetries can be presented more concretely: any configuration $x \in \mathcal{C}(\tilde{E})$ generates a bijection:

$$\theta_x = \{(l_E \tilde{e}, r_E \tilde{e}) \mid \tilde{e} \in x\}$$

This \tilde{E} gives a collection of such isomorphisms. But actually symmetry on event structures can be axiomatized directly in terms of the isomorphisms generated. So although we will occasionally refer to symmetry presented as a span as above, we will often rely on the following equivalent presentation.

1) *Isomorphism families*: An **isomorphism family** on an event structure E is a set of bijections $\theta : x \cong y$ where $x, y \in \mathcal{C}(E)$. For θ in the isomorphism family, we write $x \cong_{\theta} y$. This set of bijections has to be closed under certain operations:

- (1) *Identity*: for all $x \in \mathcal{C}(E)$, $x \cong_{\text{id}_x} x$.
- (2) *Inverse*: if $x \cong_{\theta} y$ then $y \cong_{\theta^{-1}} x$.
- (3) *Composition*: if $x \cong_{\theta} y$ then $y \cong_{\phi} z$ then $x \cong_{\theta \circ \phi} z$.
- (4) *Restriction*: if $x \cong_{\theta} y$ and $x' \subseteq x$ with $x' \in \mathcal{C}(E)$, there is a (necessarily unique) $\theta' \subseteq \theta$ and $y' \in \mathcal{C}(E)$ such that $x' \cong_{\theta'} y'$.
- (5) *Extension*: if $x \cong_{\theta} y$ and $x \subseteq x' \in \mathcal{C}(E)$, there exists a (not necessarily unique) $\theta \subseteq \theta'$ and $y' \in \mathcal{C}(E)$ such that $x' \cong_{\theta'} y'$.

Conditions (1), (2) and (3) express that isomorphisms equip the set of configurations with a groupoid structure. Conditions

(4) and (5) express that isomorphisms behave like a history-preserving bisimulation. The notion extends directly in the presence of other data such as polarity or Question/Answer labelling, which should be preserved by isomorphisms. An event structure with symmetry (*resp.* and polarity) is called an *ess* (*resp.* *essp*) for short. Event structures with a symmetry/an isomorphism family are denoted by $\mathcal{E}, \mathcal{F}, \mathcal{A}, \mathcal{B}, \dots$

We will be particularly interested in *essps* with a unique choice of extensions for positive events: an *essp* \mathcal{A} is **thin** iff for all $x \cong_{\mathcal{A}}^{\theta} y$, if θ has positive extensions to valid isomorphisms $x_1 \cong_{\mathcal{A}}^{\theta_1} y_1$ and $x_2 \cong_{\mathcal{A}}^{\theta_2} y_2$ with $\theta \sqsubseteq^+ \theta_1, \theta \sqsubseteq^+ \theta_2$, if $x_1 \cup x_2 \in \mathcal{C}(\mathcal{A})$ then $\theta_1 \cup \theta_2$ is valid as well. In other words, the left projection of the symmetry reflects positive compatibility.

2) *Maps and symmetry*: For \mathcal{E} and \mathcal{F} *ess* and $f : E \rightarrow F$ a map between the underlying event structures, we can apply f to bijections $\theta : x \cong y$ between configurations of E using:

$$f\theta = \{(f e_1, f e_2) \mid (e_1, e_2) \in \theta\}$$

We say that $f : E \rightarrow F$ **preserves symmetry** (also written $f : \mathcal{E} \rightarrow \mathcal{F}$) iff for all $x \cong_{\mathcal{E}}^{\theta} y$, $f x \cong_{\mathcal{F}}^{f\theta} f y$. There is a category \mathcal{ESS} of *ess* and maps preserving symmetry, and a category \mathcal{ESSP} in the presence of polarities. If $f, g : \mathcal{E} \rightarrow \mathcal{F}$ are parallel maps preserving symmetry we will say that they are **symmetric**, written $f \sim g$, whenever for all $x \in \mathcal{C}(E)$, we have:

$$f x \cong_{\mathcal{F}}^{\{(f e, g e) \mid e \in x\}} g x$$

3) *Games with symmetry*: Instead of playing directly on arenas, our strategies will play on the derived *games with symmetry*. A **concurrent game with symmetry** (cgs for short) is a tuple $(\mathcal{A}, \leq, \text{pol}, \lambda)$ satisfying the conditions of an arena *except negativity*, and equipped with an isomorphism family additionally assumed to be **race-free**: for all $x \cong_{\theta} y$ with a positive $\theta \cup \{(s_1^+, s_2^+)\}$ and a negative extensions $\theta \cup \{(s_3^-, s_4^-)\}$ to valid isomorphisms, then $\theta \cup \{(s_1, s_2), (s_3, s_4)\}$ is in the isomorphism family as well. Moreover, we assume that \mathcal{A} has two sub-*essp* \mathcal{A}_+ and \mathcal{A}_- on events A , such that \mathcal{A}_+ (*resp.* \mathcal{A}_-) is thin, and closed under all negative extensions of the isomorphism in \mathcal{A} (*resp.* \mathcal{A}_-). Anticipating the next subsection, this exactly means that the identity maps $\text{id}_{\mathcal{A}} : \mathcal{A}_+ \rightarrow \mathcal{A}$ and $\text{id}_{\mathcal{A}}^{\perp} : \mathcal{A}_- \rightarrow \mathcal{A}^{\perp}$ are themselves \sim -strategies. We say that \mathcal{A} and \mathcal{A}^{\perp} have **receptive thin sub-symmetries**.

Essps support operations of simple parallel composition $\mathcal{A} \parallel \mathcal{B}$ (having as events the disjoint union $\{1\} \times A \cup \{2\} \times B$ and other components inherited) and dualization \mathcal{A}^{\perp} (having polarity reversed and other components unchanged), see [4] for details – these operations preserve the conditions for cgs.

If A is an arena, we will be interested in particular in the following game with symmetry.

Definition III.1. *From an arena A , the cgs $!A$ is defined as having components $(!A, \leq, \text{pol}, \lambda)$ as described in Subsection II-C, and isomorphism family comprising bijections $\theta : x \cong y$ such that $\text{lbl} \circ \theta = \text{lbl}$, and θ preserves and reflect causal order.*

So, θ leaves invariant the label in the arena of the move played, and preserves justification – however, it can change

copy indices. With this definition, $!A$ is a cgs: its negative and positive sub-essps respectively comprise isomorphisms preserving indices of positive (*resp.* negative) events.

B. Uniform strategies

We now define uniform strategies on cgs, called \sim -strategies.

1) \sim -pre-strategies: A \sim -pre-strategy on a cgs \mathcal{A} is a map of essps:

$$\sigma : \mathcal{S} \rightarrow \mathcal{A}$$

so \mathcal{S} has to be equipped with a symmetry as well, preserved by σ . Our first condition on \sim -strategies is that they should behave as strategies in the sense of [16] at the level of *symmetries*. They should be receptive and courteous, but their action on symmetry should be receptive too: if $x \stackrel{\theta}{\cong}_{\mathcal{S}} y$, if $\sigma \theta$ extends to

$$\sigma x \cup \{\alpha\} \stackrel{\sigma \theta \cup \{(\alpha_1^-, \alpha_2^-)\}}{\cong_{\mathcal{A}}} \sigma y \cup \{\alpha'\}$$

then there are uniquely defined $s, s' \in \mathcal{S}$ such that:

$$x \cup \{s\} \stackrel{\theta \cup \{(s, s')\}}{\cong_{\mathcal{S}}} y \cup \{s'\}$$

with $\sigma s = \alpha$ and $\sigma s' = \alpha'$. We only have to require the existence, the uniqueness comes from receptivity of σ . A \sim -pre-strategy satisfying this is called **strong-receptive**.

2) *Weak equivalence*: Previously, we have seen that strategies with repetitions carry a choice of copy indices for their positive moves in $!A$. In order to satisfy the laws of cartesian closed categories (ccc) required for soundness, we need to express that strategies play the same moves *up to symmetry*.

A **weak equivalence** between two \sim -pre-strategies $\sigma : \mathcal{S} \rightarrow \mathcal{A}$ and $\tau : \mathcal{T} \rightarrow \mathcal{A}$ is given by two maps $f : \mathcal{S} \rightarrow \mathcal{T}$ and $g : \mathcal{T} \rightarrow \mathcal{S}$, making the two triangles commute up to symmetry:

$$\begin{array}{ccc} \mathcal{S} & \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \\ \xrightarrow{\sim} \end{array} & \mathcal{T} \\ \sigma \swarrow & \sim & \searrow \tau \\ & \mathcal{A} & \end{array}$$

and such that $f \circ g \sim \text{id}_{\mathcal{T}}$ and $g \circ f \sim \text{id}_{\mathcal{S}}$. Weak equivalence is a canonical notion of equivalence between \sim -pre-strategies, as it expresses isomorphism up to symmetry. We write $\sigma \simeq \tau$ if σ and τ are weakly equivalent, and $\sigma \cong \tau$ for the **isomorphism** – if all the symmetries involved are equalities.

3) \sim -strategies: Unfortunately, the equivalence relation \simeq fails to be a congruence on \sim -pre-strategies: it is not preserved by composition (introduced in the next section). Proving that \simeq is a congruence intuitively requires one to canonically transport configurations of strategies through isomorphisms in the game. As it is, symmetry ensures that such a transport exists, but not that it can be canonically chosen to induce a map. This is a very subtle point, which (among other reasons) motivated the introduction of *saturated* strategies in [4]. Saturated strategies are closed under the action of the symmetry on the game, hence providing a canonical way to transport configurations according to such isomorphisms.

Here, we follow a different route and require instead the symmetry of strategies to be as minimal as possible, formalized by the notion of *thin* essps introduced before. So formally,

a \sim -strategy on game with symmetry \mathcal{A} is a courteous, strong-receptive and *thin* \sim -pre-strategy. For an arena A we can now define formally a PCF_{\parallel} -strategy $\sigma : \mathcal{S} \rightarrow !A$ to be a \sim -strategy that is deterministic, innocent and well-bracketed.

The rest of this section is devoted to the construction of a ccc of \sim -strategies including PCF_{\parallel} -strategies. First, we will organize cgs and \sim -strategies into a compact closed category Tcg/\simeq . Restricting on cgs coming from arenas, we will form a cartesian closed sub-category Cho/\simeq (for “concurrent Hyland-Ong games”) of Tcg/\simeq , supporting the interpretation of PCF .

In Section IV we will show that PCF_{\parallel} -strategies form a sub-ccc of Cho/\simeq , that is fully abstract for PCF .

C. A category of \sim -strategies

In this subsection, we will construct a \sim -bicategory (a bicategory whose coherence laws hold up to \sim only) Tcg (for *thin concurrent games*) having as objects cgs, as morphisms \sim -strategies $\sigma : \mathcal{S} \rightarrow \mathcal{A} \parallel \mathcal{B}$ (also written $\sigma : \mathcal{A} \xrightarrow{\text{Tcg}} \mathcal{B}$, leaving \mathcal{S} implicit), and weak equivalences as 2-cells. There is also a quotient category Tcg/\simeq , however for now we refrain from quotienting: some of our later constructions will take place on concrete strategies rather than weak equivalence classes.

1) *Pullbacks and parallel interaction*: In game semantics, composition of strategies $\sigma : \mathcal{A} \xrightarrow{\text{Tcg}} \mathcal{B}$ and $\tau : \mathcal{B} \xrightarrow{\text{Tcg}} \mathcal{C}$ is usually performed by *parallel interaction*, where σ and τ are allowed to communicate on \mathcal{B} , followed by *hiding*, where only the *external* events occurring on \mathcal{A} or \mathcal{C} are retained. Traditionally, parallel interaction is formulated as an intersection of the plays that both σ and τ are prepared to play. In our games on event structures, parallel interaction is conveniently formulated as an adequate *pullback*. In general, the category of ess and maps preserving symmetry does not have all pullbacks, however it has enough of them:

Lemma III.2. *For \mathcal{A} a cgs and strong-receptive pre- \sim -strategies $\sigma : \mathcal{S} \rightarrow \mathcal{A}$ and $\tau : \mathcal{T} \rightarrow \mathcal{A}$, then temporarily forgetting about polarities there is a pullback in ESS :*

$$\begin{array}{ccc} \Pi_1 \mathcal{S} \otimes \mathcal{T} \Pi_2 & & \\ \swarrow & \vee & \searrow \\ \mathcal{S} & & \mathcal{T} \\ \sigma \searrow & & \swarrow \tau \\ & \mathcal{A} & \end{array}$$

Arbitrary maps $f : \mathcal{A} \rightarrow \mathcal{C}$ and $g : \mathcal{B} \rightarrow \mathcal{C}$ might not have a pullback, because they might not agree on possible extensions of isomorphisms in \mathcal{C} , so the extension property might fail for the natural candidate for the isomorphism family of the pullback. However in the situation described in the lemma above, strong-receptivity of σ and τ ensures that they will always find a common extension of a given isomorphism.

In this paper we will omit the construction of pullbacks, which can be found *eg.* in [20]. However we mention a key representation of their configurations and isomorphisms:

Proposition III.3. *Let \mathcal{A} be a cgs, and $\sigma : \mathcal{S} \rightarrow \mathcal{A}$ and $\tau : \mathcal{T} \rightarrow \mathcal{A}$ be strategies. Then the configurations of the*

(underlying event structure $S \otimes T$ of the) pullback $\mathcal{S} \otimes \mathcal{T}$ uniquely correspond to composite bijections:

$$x \stackrel{\sigma}{\cong} \sigma x = \tau y \stackrel{\tau}{\cong} y$$

with $x \in \mathcal{C}(S)$ and $y \in \mathcal{C}(T)$, $x \stackrel{\sigma}{\cong} \sigma x$ is the bijection induced by σ , and the composite bijection is **secured**, ie. the transitive relation generated by $(s, t) \leq (s', t')$ if $s \leq s'$ or $t \leq t'$ is a partial order – which corresponds to the order of the pullback.

Isomorphisms of $\mathcal{S} \otimes \mathcal{T}$ correspond to commuting diagrams:

$$\begin{array}{ccccc} x_1 & \stackrel{\sigma}{\cong} & \sigma x_1 & = & \tau y_1 & \stackrel{\tau}{\cong} & y_1 \\ \theta \wr \llcorner_{\mathcal{S}} & & & & & & \theta' \wr \llcorner_{\mathcal{T}} \\ x_2 & \stackrel{\sigma}{\cong} & \sigma x_2 & = & \tau y_2 & \stackrel{\tau}{\cong} & y_2 \end{array}$$

Interestingly, the symmetry $\widetilde{S \otimes T}$ is the pullback $\widetilde{S} \otimes \widetilde{T}$.

2) *Projection and hiding*: Now that we have a notion of parallel interaction, we need a way to express *hiding*. This is done via the notion of *projection*. If E is an event structure and $V \subseteq E$ is a subset of the events of E , then there is a new event structure $E \downarrow V$ whose events are those in V , and whose causal order and conflict are inherited from E [16].

If \mathcal{E} is an ess, then $V \subseteq E$ is **closed under symmetry** when for all $v \in V$, for all $x \stackrel{\theta}{\cong}_{\mathcal{E}} y$ such that $v \in x$, we have $\theta v \in V$ as well. In that case, the set comprising bijections:

$$x \cap V \stackrel{\theta \cap V^2}{\cong} y \cap V$$

for all $x \stackrel{\theta}{\cong}_{\mathcal{E}} y$ is an isomorphism family, generalizing the notion of projection in the presence of symmetry.

3) *Composition*: We are finally define the composition of \sim -strategies. Take $\sigma : \mathcal{S} \rightarrow \mathcal{A}^\perp \parallel \mathcal{B}$ and $\tau : \mathcal{T} \rightarrow \mathcal{B}^\perp \parallel \mathcal{C}$ strong-receptive \sim -pre-strategies. We follow [16] and form:

$$\begin{array}{ccc} & (\mathcal{S} \parallel \mathcal{C}) \otimes (\mathcal{A} \parallel \mathcal{T}) & \\ & \swarrow \Pi_1 \quad \vee \quad \searrow \Pi_2 & \\ \mathcal{S} \parallel \mathcal{C} & & \mathcal{A} \parallel \mathcal{T} \\ & \swarrow \sigma \parallel \mathcal{C} \quad \mathcal{A} \parallel \mathcal{B} \parallel \mathcal{C} \quad \swarrow \mathcal{A} \parallel \tau & \end{array}$$

which exists by Lemma III.2. We write $\tau \otimes \sigma = (\sigma \parallel \mathcal{C}) \circ \Pi_1$. Then, we need to hide: we do that by forming

$$V = \{v \in (\mathcal{S} \parallel \mathcal{C}) \otimes (\mathcal{A} \parallel \mathcal{T}) \mid \forall b \in \mathcal{B}, (\tau \otimes \sigma)v \neq (2, b)\}$$

so V contains the events of the pullback that map either to \mathcal{A} or \mathcal{C} , in other words which are *external*. The set V is closed under symmetry, so we can form:

$$\mathcal{T} \odot \mathcal{S} = (\mathcal{S} \parallel \mathcal{C}) \otimes (\mathcal{A} \parallel \mathcal{T}) \downarrow V$$

then $\tau \otimes \sigma : (\mathcal{S} \parallel \mathcal{C}) \otimes (\mathcal{A} \parallel \mathcal{T}) \rightarrow \mathcal{A} \parallel \mathcal{B} \parallel \mathcal{C}$ restricts to

$$\tau \odot \sigma : \mathcal{T} \odot \mathcal{S} \rightarrow \mathcal{A} \parallel \mathcal{C}$$

and reinstating polarities, we obtain $\tau \odot \sigma : \mathcal{T} \odot \mathcal{S} \rightarrow \mathcal{A}^\perp \parallel \mathcal{C}$.

Interestingly, from this definition we have the isomorphism $\widetilde{\mathcal{T} \odot \mathcal{S}} \cong \widetilde{\mathcal{T}} \odot \widetilde{\mathcal{S}}$; in fact this could even be used as a definition of the symmetry of $\mathcal{T} \odot \mathcal{S}$.

4) *Copycat*: If \mathcal{A} is an esp, there is a new esp $\mathbb{C}\mathcal{A}$ having:

- Events, those of $\mathcal{A}^\perp \parallel \mathcal{A}$,
- Causal order, the transitive closure of:

$$\begin{aligned} & \{((1, a), (1, a')) \mid a \leq a'\} \cup \{((2, b), (2, b')) \mid b \leq b'\} \\ & \cup \{((1, a), (2, a)) \mid \text{pol}(a) = +\} \\ & \cup \{((2, a), (1, a)) \mid \text{pol}(a) = -\} \end{aligned}$$

- Consistency, that of $\mathcal{A}^\perp \parallel \mathcal{A}$.

For \mathcal{A} a game with symmetry, we need to equip $\mathbb{C}\mathcal{A}$ with a symmetry. There is a canonical choice: indeed as remarked in [21], for $x, y \in \mathcal{C}(\mathcal{A})$ we have $x \parallel y \in \mathcal{C}(\mathbb{C}\mathcal{A})$ iff we have:

$$x \supseteq^+ \subseteq^- y$$

Isomorphisms on $\mathbb{C}\mathcal{A}$ are then defined as pairs of isomorphisms $x_1 \stackrel{\theta}{\cong}_{\mathcal{A}} x_2$ and $y_1 \stackrel{\theta'}{\cong}_{\mathcal{A}} y_2$ such that the diagram

$$\begin{array}{ccc} x_1 & \supseteq^+ \subseteq^- & y_1 \\ \theta \wr \llcorner_{\mathcal{A}} & & \theta' \wr \llcorner_{\mathcal{A}} \\ x_2 & \supseteq^+ \subseteq^- & y_2 \end{array}$$

commutes. Each such diagram induces a bijection between $x_1 \parallel y_1 \in \mathcal{C}(\mathbb{C}\mathcal{A})$ and $x_2 \parallel y_2 \in \mathcal{C}(\mathbb{C}\mathcal{A})$, and the isomorphism family of $\mathbb{C}\mathcal{A}$ comprises all such bijections. For arbitrary esp \mathcal{A} , this need not be a valid isomorphism family as it can fail extension. However it is valid if the isomorphism family of \mathcal{A} is **race-free**, which is the case for cgs. In terms of symmetry, these definitions imply the isomorphism $\widetilde{\mathbb{C}\mathcal{A}} \cong \mathbb{C}\widetilde{\mathcal{A}}$.

Thus from a cgs \mathcal{A} we obtain a \sim -pre-strategy $\mathcal{C}\mathcal{A} : \mathbb{C}\mathcal{A} \rightarrow \mathcal{A}^\perp \parallel \mathcal{A}$, which satisfies all the conditions to be a \sim -strategy. Our definitions of composition and copycat are compatible with those of [16], so for the same reason as there we know that there is a bicategory having arenas as objects, \sim -strategies as morphisms and isomorphisms preserving the projection to the game on the nose as 2-cells. However, we have argued before that we need to relate strategies *via* weak equivalences rather than isomorphisms, and the universal property of pullback does not ensure that weak equivalence is preserved under composition. Proving this, and setting up the framework so that it is true (while avoiding the saturation of [4]) was a significant part of the endeavour.

5) *Bipullbacks and preservation of weak equivalence*:

Preservation of weak equivalence by composition relies on the observation that the pullbacks involved in the composition also satisfy a weaker universal property up to symmetry:

Lemma III.4. *If $\sigma : \mathcal{S} \rightarrow \mathcal{A}$ and $\tau : \mathcal{T} \rightarrow \mathcal{A}^\perp$ are strong-receptive, courteous pre- \sim -strategies with receptive thin sub-symmetries, then their pullback:*

$$\begin{array}{ccc} & \mathcal{S} \otimes \mathcal{T} & \\ & \swarrow \Pi_1 \quad \vee \quad \searrow \Pi_2 & \\ \mathcal{S} & & \mathcal{T} \\ & \swarrow \sigma \quad \mathcal{A} \quad \swarrow \tau & \end{array}$$

also satisfies the universal property of bipullbacks: for all $f : \mathcal{X} \rightarrow \mathcal{S}$ and $g : \mathcal{X} \rightarrow \mathcal{T}$ making the outer square commute

up to symmetry (ie. $\sigma \circ f \sim \tau \circ g$), there exists $h : \mathcal{X} \rightarrow \mathcal{S} \otimes \mathcal{T}$, unique up to symmetry, such that $\Pi_1 \circ h \sim f$ and $\Pi_2 \circ h \sim g$.

The hypotheses of the lemma provide us, for each $x \in \mathcal{C}(X)$, with configurations $fx \in \mathcal{C}(S)$ and $gx \in \mathcal{C}(T)$ and an isomorphism $\sigma(fx) \cong_A^{\theta_x} \tau(gx)$. Roughly, the proof of the lemma constructs interactively y_S isomorphic to fx and y_T isomorphic to gx such that $\sigma y_S = \tau y_T$, so that the pullback property can be applied. Each extension of x is positive for either S or T : accordingly S or T has a canonical extension granted by its thin symmetry, to which the other strategy has to be receptive by strong-receptivity.

From this weaker universal property, it follows that weak equivalence is preserved by composition – note that applying this on the composition pullback uses the fact that cgs have receptive thin sub-symmetries. The \sim -strategies are stable under composition, and comprise cpycat. Similarly to [4], we have that Tcg is a \sim -bicategory satisfying the laws of a compact closed category up to isomorphism and weak equivalence.

D. A ccc of concurrent Hyland-Ong games

We now describe a sub- \sim -category Cho of Tcg , that gives a concurrent generalization of standard HO games.

Objects of Cho are negative arenas, ie. arenas whose minimal events are negative. A Cho-strategy from arena A to arena B (written $\sigma : A \xrightarrow{\text{Cho}} B$) is a negative \sim -strategy $\sigma : S \rightarrow !A^\perp \parallel !B$ (ie. S is negative), which is **single-threaded**:

- (1) For any $s \in S$, $[s]$ has exactly one minimal event,
- (2) For any $x \in \mathcal{C}(S)$ such that $x \xrightarrow{s_1} c, x \xrightarrow{s_2} c$ and $x \cup \{s_1, s_2\} \notin \mathcal{C}(A)$, then $[s_1] \cap [s_2] \neq \emptyset$.

Both negativity and single-threadedness are stable under composition. Negativity ensures that σ does not start playing on A , and single-threadedness will ensure that the strategy satisfies surjective pairing. Finally, 2-cells are weak equivalences.

We mention a convenient operation for constructing strategies. Take a Cho-strategy $\sigma : S \rightarrow !B^\perp \parallel !C$ and a strong-receptive, courteous map $f : !B^\perp \rightarrow !A^\perp$. Then, the composite:

$$(f \parallel !C) \circ \sigma : S \rightarrow !A^\perp \parallel !C$$

is still a Cho-strategy. Moreover if $A = B$ and $f \sim \text{id}_{!A}$, it follows that the obtained map is weakly equivalent to σ .

1) *Cartesian structure*: Since morphisms in Cho are negative strategies and arenas are negative, it follows that the empty arena 1 is terminal. From two arenas A and B , one can form the product arena $A \times B$; it is simply defined as $A \parallel B$. Firstly there are projections, obtained as:

$$\text{i}_A = (\text{i}_A \parallel !A) \circ \text{c}_{!A} : \mathbb{C}_{!A} \rightarrow !(A \times B)^\perp \parallel !A$$

where $\text{i}_A : !A^\perp \rightarrow !(A \times B)^\perp$ is the injection map, which is strong-receptive (since A, B are negative) and courteous. The other projection $\text{w}_B : A \times B \xrightarrow{\text{Cho}} B$ is defined similarly.

From $\sigma : S \rightarrow !A^\perp \parallel !B$ and $\tau : T \rightarrow !A^\perp \parallel !C$, we define their pairing $\langle \sigma, \tau \rangle : A \xrightarrow{\text{Cho}} B \times C$. For that, we need to make sure that σ and τ reach a disjoint image in $!A^\perp$. Take A an arena, and $\iota : \omega \rightarrow \omega$ an injection. Then ι induces a map

$\iota : !A \rightarrow !A$. The extended ι applies the injection to the copy indices of all minimal events, leaving the rest unchanged. More precisely, for $\alpha \in !A$, $\iota(\alpha)$ is an index function α' defined as:

$$\begin{aligned} \alpha' &: [\text{lbl } \alpha] &\rightarrow & \omega \\ a & &\mapsto & \iota(\alpha(a)) \quad \text{if } a \in \min(A) \\ a & &\mapsto & \alpha(a) \quad \text{otherwise} \end{aligned}$$

Note that for any such ι , the map $\iota^\perp : !A^\perp \rightarrow !A^\perp$ is strong-receptive and courteous, and we always have $\iota^\perp \sim \text{id}_{!A^\perp}$, by definition of the symmetry on $!A$. Take two such injections that have a disjoint codomain, e.g. $\iota_e(n) = 2n$ and $\iota_o(n) = 2n + 1$. Then, by the observation above we have Cho-strategies:

$$(\iota_e^\perp \parallel !B) \circ \sigma : A \xrightarrow{\text{Cho}} B \quad (\iota_o^\perp \parallel !C) \circ \tau : A \xrightarrow{\text{Cho}} C$$

with disjoint images in $!A^\perp$. So by co-pairing (and injection of $!B$ and $!C$ into $!(B \times C)$) we obtain:

$$\langle \sigma, \tau \rangle : S \parallel T \rightarrow !A^\perp \parallel !(B \times C)$$

which is a Cho-strategy. It follows (by a reasoning similar to that for the neutrality of cpycat) that $\text{w}_B \circ \langle \sigma, \tau \rangle$ is isomorphic to $(\iota_e^\perp \parallel !B) \circ \sigma$, which is weakly equivalent to σ since $\iota_e \sim \text{id}_{!A}$ – the same holds for $\text{w}_C \circ \langle \sigma, \tau \rangle \simeq \tau$.

Finally, any $\sigma : A \xrightarrow{\text{Cho}} B \times C$ satisfies surjective pairing:

$$\sigma \simeq \langle \text{w}_B \circ \sigma, \text{w}_C \circ \sigma \rangle$$

This relies on the fact that Cho-strategies are single-threaded. Indeed, for $\sigma : S \rightarrow !A^\perp \parallel !(B \times C)$ single-threaded, $x \in \mathcal{C}(S)$ can be decomposed as $x = x_B \cup x_C$, where $x_B, x_C \in \mathcal{C}(S)$ are disjoint (by (1)) and only reach respectively B and C . Likewise, two configurations $x_B \in \mathcal{C}(S), x_C \in \mathcal{C}(S)$ reaching only respectively B and C are compatible by (2) – the proof of the weak equivalence above follows these lines.

From all the above, it follows that Cho is cartesian up to weak equivalence. We now check that it is cartesian closed.

2) *Exponentials*: A morphism from $A \times B$ to C in Cho is a map: $\sigma : S \rightarrow !(A \times B)^\perp \parallel !C$. Up to isomorphism, its codomain is $!A^\perp \parallel !(B^\perp \parallel C)$, which suggests $B^\perp \parallel C$ as an exponential object. However, $B^\perp \parallel C$ is not a valid arena, since it fails negativity – and negativity is essential to ensure that the empty arena 1 is a terminal object.

The standard arena construction $A \Rightarrow B$ introduced in Subsection II-C solves this issue by setting minimal events of B as dependencies for events in A . Showing that this defines an exponential of A and B relies on the following lemma.

Lemma III.5. *There is a bijection Φ preserving and reflecting weak equivalence between Cho-strategies $\sigma : C \xrightarrow{\text{Cho}} A \Rightarrow B$, and negative single-threaded \sim -strategies $\sigma' : !C \xrightarrow{\text{Tcg}} !A^\perp \parallel !B$. Moreover, for all $\tau : D \xrightarrow{\text{Cho}} C$, $\Phi(\sigma \circ \tau) \simeq \Phi(\sigma) \circ \tau$.*

Proof. First, we note that for arenas A and B there is a strong-receptive, courteous map of essp :

$$\chi_{A,B} : !(A \Rightarrow B) \rightarrow !A^\perp \parallel !B$$

Events $\beta : [(2, b)] \rightarrow \omega$ are kept unchanged. However, we have multiple copies of A on the left hand side and

only one on the right hand side, so we need to reindex minimal events of A to avoid collisions. Therefore we send events $\alpha : [(1, (b, a))] \rightarrow \omega$ with a minimal to $(1, \alpha')$ with $\alpha'(a) = \langle \# b, \alpha((2, b)), \alpha((1, (b, a))) \rangle$ where $\# : B \rightarrow \omega$ is given by countability of B , and $\langle -, - \rangle : \omega^3 \rightarrow \omega$ is any injection. If a is non-minimal, its copy index is unchanged.

For $\sigma : \mathcal{S} \rightarrow !A^\perp \parallel !(B \Rightarrow C)$, $\Phi(\sigma)$ is $(!A^\perp \parallel \chi_{B,C}) \circ \sigma$. For $\sigma : \mathcal{S} \rightarrow !A^\perp \parallel !(B^\perp \parallel !C)$, by (I) of single-threadedness any $s \in \mathcal{S}$ has a unique minimal dependency mapping to a minimal event of $!C$. Thus σ factors uniquely through $!A^\perp \parallel \chi_{B,C}$ up to symmetry, from which the lemma follows. \square

From the lemma above, the cartesian structure (up to \simeq) of Cho and the compact closed structure (up to \simeq) of Tcg, it follows that Cho is cartesian closed (up to \simeq).

E. Recursion

Usually, the interpretation of the fixpoint combinator Y is obtained by showing that the ccc of strategies is enriched over a category of sufficiently complete partial orders. Once this is established, the recursive equation for Y can be solved by computing the right least upper bound. Here however (as in AJM games [1]) the ccc Cho/ \simeq is a quotient, and it is not clear that the natural induced order on equivalence classes has the adequate completeness properties. Instead we use an ordering on concrete strategies rather than equivalence classes.

Definition III.6. Let $\sigma : \mathcal{S} \rightarrow \mathcal{A}$ and $\tau : \mathcal{T} \rightarrow \mathcal{A}$ be negative single-threaded \sim -strategies on a cgs \mathcal{A} . Then $\sigma \trianglelefteq \tau$ iff we have the inclusion map $\mathcal{S} \hookrightarrow \mathcal{T}$ with all data in \mathcal{S} coinciding with the restriction of that in \mathcal{T} , and for all $s \in \mathcal{S}$, $\sigma s = \tau s$.

The \sim -strategies on \mathcal{A} ordered by \trianglelefteq form a directed complete partial order (dcpo). It is not pointed though – it does not have a least element: indeed there is one minimal \sim -strategy on \mathcal{A} for each renaming of its minimal negative events. We call such a minimal \sim -strategy **empty**. For each \mathcal{A} we will consider one empty \sim -strategy in particular, written $\perp_{\mathcal{A}} : \min^-(\mathcal{A}) \rightarrow \mathcal{A}$, that acts as the identity on the ess $\min^-(\mathcal{A})$ having as events the minimal negative events of \mathcal{A} , with isomorphism family closed by receptivity. Not every \sim -strategy $\sigma : \mathcal{S} \rightarrow \mathcal{A}$ is above $\perp_{\mathcal{A}}$, however there is always $\sigma^\dagger : \mathcal{S} \rightarrow \mathcal{A}$ such that $\sigma^\dagger \simeq \sigma$ and $\perp_{\mathcal{A}} \trianglelefteq \sigma^\dagger$, obtained by renaming minimal events of \mathcal{S} .

Let us write $\mathcal{D}_{\mathcal{A}}$ for the pointed dcpo of \sim -strategies above $\perp_{\mathcal{A}}$. For an arena A we can now define the following operation, using the ccc combinators of Cho:

$$\begin{aligned} F & : \mathcal{D}_{!(A \rightarrow A)^\perp \parallel !A} \rightarrow \mathcal{D}_{!(A \rightarrow A)^\perp \parallel !A} \\ \sigma & \mapsto (ev_{A,A} \odot \langle \mathcal{C}_{!(A \rightarrow A)}, \sigma \rangle)^\dagger \end{aligned}$$

Then F is a continuous function, and has a least upper bound $Y = \nu F : A \Rightarrow A \xrightarrow{\text{Cho}} A$, which is easily shown using ccc laws to be a fixpoint combinator up to weak equivalence.

F. Reduced form for PCF $_{\parallel}$ -strategies

Describing Cho-strategies for basic combinators of PCF can be challenging, as one has to give the isomorphism family. However, we show here that in the presence of visibility, innocence and determinism, full Cho-strategies can be recovered

from the more compact (symmetry-free, Opponent replication free) notion of *reduced* PCF $_{\parallel}$ -strategies used in Section II – a situation analogous to the presentation of sequential innocent strategies as sets of P-views (compact, not compositional) or as sets of plays (not compact, compositional).

First of all, let us take a PCF $_{\parallel}$ -strategy $\sigma : \mathcal{S} \rightarrow !A$. Its **reduced form**, written \mathcal{S}_{rf} , is the sub-event structure of \mathcal{S} comprising all events $s \in \mathcal{S}$ such that for all $t^- \leq s$, $\text{ind}(\sigma) = 0$. In other words, \mathcal{S}_{rf} is the restriction of \mathcal{S} where Opponent moves always have a copy index of 0 – there is no Opponent duplication. The map $\sigma : \mathcal{S} \rightarrow !A$ also induces a map of esp $\sigma_{\text{rf}} : \mathcal{S}_{\text{rf}} \rightarrow !^+A$. The event structure \mathcal{S}_{rf} inherits a symmetry from \mathcal{S} ; however then this symmetry is trivial. Indeed if $x, y \in \mathcal{C}(\mathcal{S}_{\text{rf}})$ are such that $x \stackrel{\theta}{\simeq}_{\mathcal{S}} y$, then from σ thin it follows that $x = y$ and $\theta = \text{id}_x$. The symmetry on \mathcal{S}_{rf} is irrelevant, so σ_{rf} is a reduced PCF $_{\parallel}$ -strategy in the sense of Subsection II-D.

From reduced PCF $_{\parallel}$ -strategies we can build full strategies.

Lemma III.7. For a reduced PCF $_{\parallel}$ -strategy $\sigma : \mathcal{S} \rightarrow !^+A$, there is a PCF $_{\parallel}$ -strategy $\bar{\sigma} : \bar{\mathcal{S}} \rightarrow !A$ such that $\bar{\sigma}_{\text{rf}} \simeq \sigma$.

Proof. The underlying event structure for $\bar{\mathcal{S}}$ has for events indexing functions $\alpha : [s] \rightarrow \omega$ such that for all $t^+ \leq s$, $\alpha(t) = 0$. Those are causally ordered as for $!A$, and all finite sets are consistent. The symmetry is designed as for $!A$.

The map $\bar{\sigma}$ sends $(\alpha^- : [s^-] \rightarrow \omega) \in \bar{\mathcal{S}}$ to the event of $!A$ with $\text{lbl}(\sigma s)$ as label and $\text{ind} \alpha$ as copy index. For $\alpha^+ : [s^+] \rightarrow \omega$, $\bar{\sigma} \alpha^+$ has label $\text{lbl}(\sigma s)$ and as index an injective function of the indices of σs and the negative dependencies of s . We get a PCF $_{\parallel}$ -strategy satisfying the required equation. \square

In fact, there is a bijection (up to weak equivalence) between PCF $_{\parallel}$ -strategies and reduced forms thereof. We define a PCF $_{\parallel}$ -strategy to be **finite** whenever their reduced form is.

Any PCF $_{\parallel}$ -strategy $\sigma : \mathcal{S} \rightarrow !A$ is also a Cho-strategy: it is necessarily negative and is single-threaded by innocence.

G. Interpretation of PCF

Any term $x_1 : A_1, \dots, x_n : A_n \vdash M : A$ of PCF is interpreted as a Cho-strategy $\llbracket M \rrbracket : \mathcal{S} \rightarrow !\prod_{1 \leq i \leq n} A_i^\perp \parallel !A$. The interpretation of PCF follows from the ccc structure, and from the strategies for the basic combinators of PCF described through their reduced forms in Figure 2 (op stands for succ or pred). In these diagrams we omit the copy index whenever it is 0. Missing from this figure is the reduced form of the strategy for if, which appears in Subsection II-C.

Theorem III.8. For all $\vdash M : \mathbb{X}$, $M \Downarrow$ iff $\llbracket M \rrbracket$ is non-empty.

Proof. From left to right (soundness), it follows from the ccc laws up to weak equivalence and the fact that the strategies of Figure 2 obey the required elementary equations. From right to left (adequacy), it follows from the same properties, with a straightforward use of logical relations. \square

IV. FULL ABSTRACTION FOR PCF

We start this final section by the following proposition:

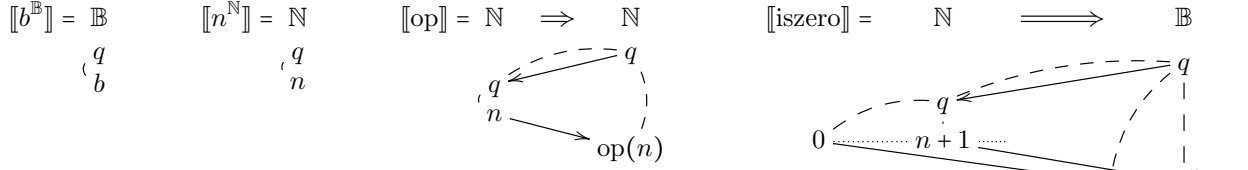


Fig. 2. Interpretation of the basic combinators of PCF

Proposition IV.1. *There is a sub- \sim -bicategory PcfPar of Cho with arenas as objects, negative PCF_{\parallel} -strategies as morphisms, and weak equivalences as 2-cells. Moreover, PcfPar satisfies the laws of a ccc up to weak equivalence.*

Proof. The difficult part is checking that composition preserves our conditions on strategies. Visibility and determinism are both stable under composition. Assuming visibility, so are innocence and well-bracketing. \square

Now, we prove that the extensional quotient of PcfPar is fully abstract for PCF. From now on we will often write $\sigma : \mathcal{A}$ instead of $\sigma : \mathcal{S} \rightarrow \mathcal{A}$ and leave \mathcal{S} implicit.

We write $M \simeq_{\text{obs}} N$ for observational equivalence between terms, defined as usual [10]. We also define it on PCF_{\parallel} -strategies by defining a **test** on arena A as any negative PCF_{\parallel} -strategy $\alpha : !A^{\perp} \parallel !\mathbb{X}$. For $\sigma : !A$ and $\sigma' : !A$ we say that σ and σ' are observationally equivalent, written $\sigma \simeq_{\text{obs}} \sigma'$, whenever for all $\alpha : !A^{\perp} \parallel !\mathbb{X}$, $\alpha \circ \sigma \simeq \alpha \circ \sigma'$. For standard reasons, \simeq_{obs} is preserved by all our constructions on strategies. To prove full abstraction, we prove a finite definability result up to \simeq_{obs} .

Take a type $A = A_1 \rightarrow \dots \rightarrow A_n \rightarrow \mathbb{X}$ of PCF, with $A_i = A_{i,1} \rightarrow \dots \rightarrow A_{i,p_i} \rightarrow \mathbb{X}_i$. An alternative presentation of the standard finite definability argument for PCF follows from the remark that any innocent $\sigma : A$ can be decomposed into a **flow substrategy** $\sigma_{\text{flow}} : \mathbb{X}_1 \rightarrow \dots \mathbb{X}_n \rightarrow \mathbb{X}$ (comprising the plays of σ where Opponent only plays answers), and for each Player question q in σ_{flow} in \mathbb{X}_i , the data of **argument sub-strategies** $\sigma_{q,j} : A_1 \rightarrow \dots \rightarrow A_n \rightarrow A_{i,j}$ ($1 \leq j \leq p_i$) detailing Player's behaviour if Opponent was to interrogate the j -th argument of A_i . To PCF-define σ , it suffices to PCF-define σ_{flow} and to inductively PCF-define the argument sub-strategies. Our finite definability process will follow these lines: we will first address the first-order case, and generalize to higher-order using a decomposition argument.

1) *First-order case:* We remark that on **first-order types** (ie. types of the form $\mathbb{X}_1 \rightarrow \dots \rightarrow \mathbb{X}_n \rightarrow \mathbb{X}$), PCF_{\parallel} -strategies denote continuous functions between the corresponding Scott domains, and that the functions definable by PCF_{\parallel} -strategies are exactly the PCF-definable functions (although they are not computed sequentially). We assume familiarity with the interpretation of PCF in Scott domains [15]. For a PCF type A , we write $\llbracket A \rrbracket_{\mathcal{D}}$ for the corresponding domain.

Firstly, any PCF_{\parallel} -strategy on a ground type must be (by determinism and well-bracketing) either empty or weakly equivalent to the interpretation of a constant. We write \uparrow

$\sigma \in \llbracket \mathbb{X} \rrbracket_{\mathcal{D}}$ for the corresponding domain element. Likewise any $d \in \llbracket \mathbb{X} \rrbracket_{\mathcal{D}}$ corresponds to a strategy $\downarrow d : !\mathbb{X}$ defined by the interpretation of the corresponding PCF combinator. We generalise \downarrow to first-order types by, for $\sigma : \prod_{i \leq n} \mathbb{X}_i \xrightarrow{\text{PcfPar}} \mathbb{X}$:

$$\downarrow \sigma : \prod_{i \leq n} \llbracket \mathbb{X}_i \rrbracket_{\mathcal{D}} \rightarrow \llbracket \mathbb{X} \rrbracket_{\mathcal{D}} \\ \langle x_i \mid i \leq n \rangle \mapsto \downarrow (\sigma \circ (\uparrow x_i \mid i \leq n))$$

The following lemma is a variant of the familiar “linear tests suffice” argument in innocent game semantics.

Lemma IV.2. *For $\sigma, \tau : \prod_{i \leq n} \mathbb{X}_i \xrightarrow{\text{PcfPar}} \mathbb{X}$, $\sigma \simeq_{\text{obs}} \tau$ iff $\downarrow \sigma = \downarrow \tau$.*

A **ground term** is a term $x_1 : \mathbb{X}_1, \dots, x_n : \mathbb{X}_n \vdash M : \mathbb{X}$ of PCF extended with a constant $\perp_{\mathbb{X}} : \mathbb{X}$ for divergence, which contains no subterm of functional type. For a ground term $\Gamma \vdash M : \mathbb{X}$, it is direct to prove by induction that $\downarrow \llbracket M \rrbracket = \llbracket M \rrbracket_{\mathcal{D}}$.

Finally, we prove finite definability on first-order types.

Lemma IV.3. *If $\sigma : \prod_{i \leq n} \mathbb{X}_i \xrightarrow{\text{PcfPar}} \mathbb{X}$ is finite negative, there is a ground $x_1 : \mathbb{X}_1, \dots, x_n : \mathbb{X}_n \vdash M_{\sigma} : \mathbb{X}$ s.t. $\llbracket M_{\sigma} \rrbracket \simeq_{\text{obs}} \sigma$.*

Proof. Exploiting determinism of σ , we construct by induction a ground term M_{σ} performing the same queries as σ in some sequential order, such that $\llbracket M_{\sigma} \rrbracket_{\mathcal{D}} = \downarrow \sigma$. By Lemmas IV.2 and the remark above, it follows that $\llbracket M_{\sigma} \rrbracket \simeq_{\text{obs}} \sigma$. \square

2) *Higher-order and full abstraction:* Let $\sigma : \prod_{i \leq n} A_i \xrightarrow{\text{PcfPar}} \mathbb{X}$ be a finite PCF_{\parallel} -strategy. We first extract:

$$\mathcal{S}_{\text{flow}} = \{s \in S \mid \sigma([s]) \subseteq !(\prod_{1 \leq i \leq n} \mathbb{X}_i)^{\perp} \parallel !\mathbb{X}\}$$

It inherits all the components of an ess from \mathcal{S} , and gives a negative PCF_{\parallel} -strategy from $\prod_{1 \leq i \leq n} \mathbb{X}_i$ to \mathbb{X} .

We call events $q^{(+, \text{Qu})} \in \mathcal{S}_{\text{flow}}$ whose negative dependencies all have copy index 0 the **primary questions** of \mathcal{S} . Each such q maps to an initial question of some A_i , we write Q_i for the set of primary questions mapping to A_i . From there, Opponent can ask for an arguments of type $A_{i,j}$ via an event $q_{q,j}^k$ (for all k) immediately dependent on q . We set:

$$\mathcal{S}_{q,j} = \{s \in S \mid \exists k \in \omega, s \geq q_{q,j}^k\}$$

It inherits an essp structure from \mathcal{S} . We reassign:

$$\sigma_{q,j} : \mathcal{S}_{q,j} \rightarrow !(\prod_{1 \leq i \leq n} A_i)^{\perp} \parallel !A_{i,j}$$

and check that $\sigma_{q,j}$ is a negative PCF_{\parallel} -strategy – it is the j -th **argument sub-strategy** of q .

We now reconstruct σ from the flow and argument sub-strategies using the ccc combinators of PcfPar . We set:

$$\sigma_{\text{flow}} : \mathcal{S}_{\text{flow}} \rightarrow !(\Pi_{1 \leq i \leq n} \Pi_{q \in Q_i} \mathbb{X}_i)^{\perp} \parallel !\mathbb{X}$$

The assignment σ_{flow} follows canonically from σ and the symmetry on \mathcal{S} , which to any Player question $q^{(+, \text{Qu})} \in \mathcal{S}_{\text{flow}}$ depending on minimal event q^{i+1} , associates canonically some $q \in Q$. With this, we can finally state our decomposition result:

Proposition IV.4. For $\sigma : \Pi_{1 \leq i \leq n} A_i \xrightarrow{\text{PcfPar}} \mathbb{X}$ a PCF_{\parallel} -strategy,

$$\sigma \simeq \sigma_{\text{flow}} \odot \langle x_i \sigma_{q,1} \dots \sigma_{q,p_i} \mid 1 \leq i \leq n, q \in Q_i \rangle$$

with σ_{flow} and $\sigma_{q,j}$ PCF_{\parallel} -strategies. The notation refers to ccc combinators, with x_i the i -th projection from $\Pi_{1 \leq i \leq n} A_i$.

The proof requires a careful analysis of the right-hand side composition, and relies on all the conditions of PCF_{\parallel} -strategies. From this immediately follows by induction:

Proposition IV.5. If $\sigma : \mathcal{S} \rightarrow !\llbracket A \rrbracket$ is a finite PCF_{\parallel} -strategy, there exists $\vdash M_{\sigma} : A$ such that $\llbracket M \rrbracket \simeq_{\text{obs}} \sigma$.

From the above ingredients we obtain easily:

Theorem IV.6. $\text{PcfPar}/\simeq_{\text{obs}}$ is fully abstract for PCF , ie. for all $\Gamma \vdash M : A, \Gamma \vdash N : A, M \simeq_{\text{obs}} N$ iff $\llbracket M \rrbracket \simeq_{\text{obs}} \llbracket N \rrbracket$.

Proof. A standard argument using soundness, adequacy, finite definability, and the easy fact that if two strategies can be distinguished by a test, then a finite test suffices. \square

V. CONCLUSION

We described a ccc Cho/\simeq of concurrent non-deterministic strategies, which supports concurrent notions of visibility, well-bracketing and innocence that conservatively extends the usual ones. In this setting, we gave a parallel intensionally fully abstract interpretation of PCF .

As a first extension, we plan to give an intensional version of the fully abstract model of PCF extended with the concurrent primitive parallel-or [15]. This is trickier than it might seem: the concurrent strategy for parallel-or is non-deterministic, and linking its non-deterministic intensional behaviour to its deterministic extensional behaviour requires some care.

More generally, we plan to apply this framework to give fine-grained truly concurrent games model of higher-order effectful concurrent programming languages.

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APPENDIX

The detailed construction of concurrent games with symmetry and thin strategies and of the category Cho is available on arXiv: <http://arxiv.org/abs/1409.7542>. All proofs concerning the conditions on strategies and the full-abstraction result can be found at <http://perso.ens-lyon.fr/pierre.clairambault/pcfpar.pdf>.