

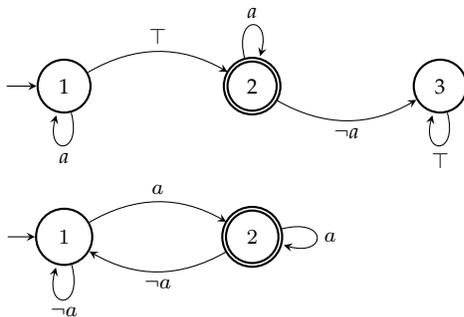
## TD4: GNBA's

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**Exercise 1** (Automata). — 1. Give NBAs recognizing the same language as the following formulas:

- $\mathbf{G}(a \rightarrow Fa)$
- $a\mathbf{W}b = \mathbf{G}a \vee (a\mathbf{U}b)$

2. For  $a\mathbf{W}b$ , build the corresponding GNBA and compare.
3. Give LTL formulas equivalent to the following automata:



**Exercise 2** (LTL versus NBAs). — Consider the language  $L = \{q_0 \dots q_n \mid a \in q_{2i}\}$  on the set of atomic formulae  $\text{At} = \{a\}$ .

1. Give a NBA recognizing  $L$
2. Prove that for any LTL-formula  $\varphi$ , there exists  $n \in \mathbb{N}$  such that  $\{a\}^{n+1} \mathcal{O} \{a\}^\omega \models \varphi$  if and only if  $\{a\}^{n+2} \mathcal{O} \{a\}^\omega \models \varphi$ .
3. Prove that NBAs are strictly more expressive than LTL formulas.
4. Show that  $[\alpha] \mapsto \mathbf{X}\mathbf{X}\alpha \wedge a$  has no greatest fixpoints. Deduce that the lattice of LTL formulae is not complete.

**Exercise 3** (Into the past). — We write  $\text{At}$  for the set of atomic propositions.

1. We wish to extend the syntax of LTL with two more modalities  $\varphi\mathbf{S}\psi$  (“since”) and  $\mathbf{P}\varphi$  (“previous”). Explain how we could extend the semantics of LTL to accommodate those new constructors.
2. Explain how to define connectors to capture the informal properties: “at some point in the past” and “always in the past”.
3. Show how to translate a formula  $\varphi$  of LTL+past operators in  $\varphi'$  where the past operators are not nested by augmenting the set of atomic formulas to  $\text{At}'$ . How to state correctness of our translation?
4. Show how to reduce a formula of LTL+past operators  $\varphi$  into  $\varphi'$  without past operators, augmenting the set of atomic propositions to  $\text{At}'$ , that is correct in the sense of the previous question.