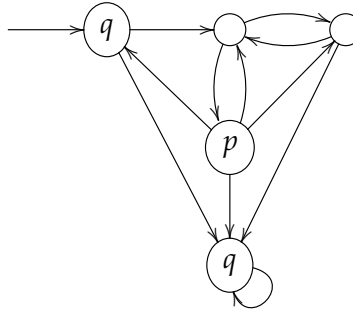


TD5: CTL

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Exercise 1 (Warmup). — Consider the following transition system \mathcal{T} ($At = \{p, q\}$):



Determine which of the following formulas are satisfied by \mathcal{T} :

 EFp AGp $AXEFp$ $EXEFp$ $E(qUXAp)$ $A(qUXAp)$

Exercise 2 (CTL equivalences). — Two formulas of CTL are equivalent when they are satisfied by the same LTSs. Show the following equivalences:

1. $AF\varphi \equiv \neg EG\neg\varphi$
2. $AF\varphi \equiv \varphi \vee AXAF\varphi$
3. $AG\varphi \equiv \varphi \wedge AXAG\varphi$

Exercise 3 (Comparing LTL and CTL). — 1. Are the LTL formula GFp and the CTL formula $AGAFp$ equivalent?

2. Are the LTL formula FGp and the CTL formula $AFAGp$ equivalent?
3. Are the LTL formula $G(p \rightarrow Fq)$ and the CTL formula $AG(p \rightarrow AFq)$ equivalent?
4. Are the LTL formula $F(p \wedge Xp)$ and the CTL formula $AF(p \wedge AXp)$ equivalent?

Exercise 4 (Difference of expressivity). — 1. Show that there exists a formula φ of LTL such that there exists no CTL formula ψ such that for all states s of all LTS, $s \models \varphi$ if and only if $s \models \psi$.

2. Show that there exists a formula φ of CTL such that there exists no LTL formula ψ such that for all states s of all LTS, $s \models \varphi$ if and only if $s \models \psi$.

Exercise 5 (LTL is not complete). — Recall the function $g([\alpha]) = [XX\alpha \wedge a]$ on $At = \{a\}$ and the language $L = \{w \in \Sigma^\omega \mid \forall i, w_{2i} \neq \emptyset\}$ which is known not to be LTL-definable (see previous TD).

Assume the lattice LTL / \equiv is complete.

1. Show $[Ga] \leq [g([Ga])]$ and $\text{Words}(Ga) \subseteq L$.
2. Let $\varphi = \bigvee_{n \in \mathbb{N}} g^k([Ga])$. Show that $\text{Words}(\varphi) = L$.
3. Conclude.